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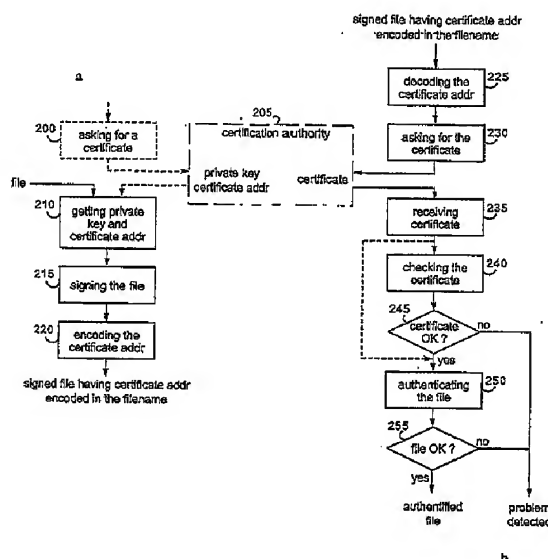
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[Continued on next page]

(54) Title: **METHOD AND SYSTEM FOR LINKING CERTIFICATES TO SIGNED FILES**



(57) Abstract: A method and systems for linking the certificate to the signed file to which it is associated is disclosed. According to the method and systems of the invention, the address or URL wherein the certificate is stored is encoded in the filename of the signed file so as to be transmitted jointly with the file. When receiving such a signed file, a first step consists in extracting and decoding the certificate address from the filename. Using the certificate address, the certificate can be accessed and checked before opening and authenticating the signed file. In a preferred embodiment, the signature of a signed file is based on the file content and on a private key while the corresponding certificate comprises at least the corresponding public key.

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METHOD AND SYSTEM FOR LINKING CERTIFICATES TO SIGNED FILES

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Field of the Invention

The present invention relates generally to network computing security and more specifically to a method and systems for linking a digital certificate to a digitally
15 signed file that can be accessed through a network so as to provide information relative to the signer identity and the validity of the signature that can be used before opening the file.

Background of the Invention

20 To improve data transmission security over computer networks and to prevent digital forgery, a digital signature is commonly used to authenticate a file i.e., to check file integrity and to authenticate signer. Such digital signature allows, for example, to control the source of a received
25 file, and to verify the file integrity. A digital signature asserts that the user corresponding to the digital signature

wrote or otherwise agreed with the contents of an electronic document or other information object to which the digital signature is appended. As with written signatures, digital signatures provide authentication of the signer's identity, acceptance of the terms stated in the signed document, proof of the integrity of the document's contents, and non repudiation (in other words, the signer cannot deny what he/she has signed). Digital signatures are generally based upon public key algorithms wherein security is provided through keys independently of the used algorithm, which may be freely published or analyzed.

A digital certificate can be considered as an attachment to a signed document, to link the identity of the signer of the document to his/her public key. A digital certificate provides a cryptographic public key that allows another party to encrypt information for the certificate's owner. A digital certificate also allows to verify that a user sending a document is who he/she claims to be, and to provide the receiver with the means to encode a reply. A certificate therefore securely identifies the owner of the public key pair, which is used to provide authentication, authorization, encryption, and non-repudiation services. A digital certificate contains the signer's public key and bears the digital signature of a Certification Authority (CA). The most widely used standard for digital certificates is X.509, Version 3, "The Directory-Authentication Framework 1988", promulgated by the International Telecommunications Union (ITU), which defines the following structure for public-key certificates:

- version field (identifying the certificate format)
- Serial Number (unique within the CA)

- Signature Algorithm (identifying the issuer's hash and digital signature algorithms used to sign the certificate)
- Issuer Name (the name of the CA)
- Period of Validity (a pair of "Not Before", and "Not After" Dates)
- Subject Name (the name of the user to whom the certificate is issued)
- Subject's Public Key field (including Algorithm name and the Public Key of the subject)
- Extensions
- Signature of CA

A certification authority is the third party that everyone trusts whose responsibility is to issue digital certificates providing the link between the signer and the signer's public key. A certification authority (CA) also keeps records about the transactions that occur using certificates it has issued. An individual wishing to sign a document applies for a digital certificate from a Certification Authority. The digital certificate is digitally signed by the issuing Certification Authority that ensures both content and source integrity. The CA makes its own public key readily available through, for example, print publicity or on the Internet. The act of digitally signing makes the certificates substantially tamperproof, and therefore further protection is not needed. The strength of protection equates directly to the strength of the algorithm and key size used in creating the issuer's digital signature (hash and digital signature algorithms).

The signature verification process checks the digital signature 'appended or attached to a document using the public verification key extracted from the digital certificate, issued by the CA, that must be also appended to or referenced in the document. Using the public key of the signer, the signature verification process recovers from the digital signature, the hash value, computed by the signer, in the file that was signed using the private key of the signer during the authentication process. To verify that the file is authentic, the receiver computes also the hash value of the document, and then compares the deciphered hash value with the real hash value, computed from the file. If both hash values are identical, the file is accepted as authentic, otherwise, the file is rejected as being corrupted or fake.

Once the digital signature of a file has been computed and the file has been signed with the digital signature for verification purposes, a digital certificate must be associated with the signed file to make possible the verification of the digital signature by the recipient.

Generally, a digital certificate used for authenticating a file is transmitted as a separate file, appended to the file it authenticates e.g., as part of a file wrapper structure, or alternatively, the certificate can be retrieved from a reference or address e.g., the URL of the certificate on the issuing CA Web Server.

Transmitting and maintaining digital certificates and signed documents as separate files e.g., the digital certificate associated to a signed document is stored in the user's workstation or in a server, presents the advantage of

supporting file authentication at any time in a simple and well understood way. However, if documents are later passed on or moved to new recipients, associated digital certificates can be lost, accidentally removed, or even intentionally removed on the way in an attempt to cheat.

Wrapping a file with delimiters and appending the digital certificate, or the URL of said certificate on the issuing CA Web Server, at the end of the signed file is convenient, since both the certificate, or the certificate address, and the signed content travel together. Conversely, the wrapper and the certificate, or the certificate address, will typically need to be removed before the file can be used. Thus, signature validation only occurs when the document is retrieved. If the document is later passed on or moved, it may be difficult to check again, since the certificate, or the certificate address, could be lost. Furthermore, the method is not compatible with standard file formats such as image, video, audio or executable files that cannot be recognized prior to authentication.

When a recipient receives an electronic document, if the digital certificate has been appended to the signed document, the recipient must perform the following tasks:

- open the electronic document;
- 5 - identify and extract, from the electronic document, the digital certificate and the digital signature portions appended to this electronic document;
- identify the address and contact the CA to check that the appended certificate is a valid certificate, using the
- 10 digital certificate content; and,
- verify the signature using the public key in the certificate.

It must be observed that if the digital certificate is appended to the received electronic document, the recipient

15 must open the document file for accessing the digital certificate required to verify the signature. Even when the certificate, instead of being appended, would be referenced e.g., as a network address or URL, in the received document, the address from which the certificate e.g., from a CA Web

20 Server or directory archive, can be accessed or retrieved, must also be appended by the sender to the signed document. Therefore, it is also required to open the received document to get said address needed for accessing the digital certificate.

25 Thus, there are security problems related to the methods described above for verifying the authenticity of received or accessed files by the recipient:

- when certificates are sent as separate files, the associated digital certificates could be lost if the
- 30 signed files are later passed on or moved to new

recipients. In such case, it is impossible to verify these signed files.

5 - when certificates, or certificates addresses, are
appended to the signed files, recipients must open and
process the received files to verify said files. Before
opening a received files, parsing the content for
locating, and retrieving, or accessing, the associated
certificate, there is no way to determine in advance,
whether the received file has been signed or not i.e.,
10 whether it is an "authenticated" file or an "impersonated"
file (a non-signed file). Likewise, it is impossible to
determine whether or not the certificate is valid i.e., if
it has been issued by a CA, if it has not been revoked,
and if the certificate date is valid.

15 It is also to be noticed that opening files for verifi-
cation represents an important security concern.

Many viruses spread on the Internet on e-mail attach-
ments distributed as "impersonated". If a received imperson-
ated file has been maliciously infected by a virus, opening
20 the infected file for the simple purpose of signature
verification almost surely may "open" the door for infecting
the receiver's computer. This is a "security hole" common to
all signature methods described above, as illustrated by
operation of the class of public-key algorithms discussed
25 herein before.

Certificates must be issued by certificate authorities.
If a certificate becomes compromised, the certificate
authority can later revoke the certificate, thus rendering
invalid all files signed after the signature's revocation
30 date. A certificate could become compromised if an

unauthorized third-party obtained the private key associated with the certificate. This private key is typically stored on the signer's computer. With the private key, an unauthorized person could essentially forge a signature. If the
5 recipient receives a file signed with a revoked certificate, it is must be discarded as invalid or fake.

Therefore, before opening a received file, it would be advisable to check:

- 10 - if the file has been signed i.e., if it contains a digital signature and a digital certificate appended or referenced;
- the issuer name i.e., the name of the CA;
- the name of the user to whom the certificate has been issued; and,
- 15 - the validity period of the certificate.

Therefore, there is a need to provide a method and systems for accessing a digital certificate from a signed file before opening said file, so as to enable the recipient of the file to determine if the received file has been
20 signed i.e., authenticated, and to check the identify of signer e.g., contacting the signer by e-mail, and the validity of the digital certificate before opening said file for signature verification.

Summary of the Invention

25 Thus, it is a broad object of the invention to remedy the shortcomings of the prior art as described here above.

It is another object of the invention to provide a method and systems adapted for enabling a recipient to check whether or not a received file is a signed file, before opening said file.

- 5 It is a further object of the invention to provide a method and systems adapted to access the digital certificate of a signer, before opening the corresponding files.

10 It is still a further object of the invention to provide a method and systems adapted for providing the identity and address of the signer of a file to the recipient of this file so as to verify signer identity, before opening a suspicious file.

15 The accomplishment of these and other related objects is achieved by a method for encoding in the filename of a signed file, an address from which the certificate required to authenticate said signed file can be accessed, said method comprising the steps of,

- encoding said address from which the certificate required to authenticate said signed file can be accessed;
 - 20 - merging said filename and said encoded address in a new filename; and,
 - renaming said signed file with said new filename,
- wherein said filename and said encoded addresses are separated by a control character,

25 and by a method for authenticating a signed file having a filename wherein an address from which the certificate required to authenticate this signed file can be accessed is encoded, said method comprising the steps of,

- extracting said encoded address;
- decoding said encoded address;
- accessing said certificate required to authenticate said signed file using said decoded address,
- 5 - authenticating said signed file using said accessed certificate.

Further advantages of the present invention will become apparent to the ones skilled in the art upon examination of the drawings and detailed description. It is intended that
10 any additional advantages be incorporated herein.

Brief Description of the Drawings

- Figure 1** , comprising figures 1a, 1b, and 1c, illustrates an example of the algorithm used for encoding, in the filename of a file, the address or URL wherein the certificate used to signed this file is stored.
- Figure 2** , comprising figures 2a and 2b, illustrates an example of the algorithm that is used to sign an electronic document and of the algorithm that is used to check the integrity and to verify the signature of a signed file, respectively.
- Figure 3** depicts an example of the environment wherein the invention can be implemented.
- Figure 4** shows an example of a signed file content wherein the filename is encoded according to the invention.
- Figure 5** , comprising figures 5a to 5f, illustrates an example of the user's interface when using the invention.

Detailed Description of the Preferred Embodiment

According to the invention, the filename of a file that is accessed locally or through a computer network is used to encode the address, or URL, from which the certificate that
5 can be used to check the integrity and to verify the signature of the file can be accessed. A lexicography is determined so as to avoid particular characters that may be forbidden by the file system, e.g., "\" with Microsoft Windows system (Windows is a Trademark of Microsoft Corpora-
10 tion), and/or to encode the addresses so as to reduce their sizes. Addresses to be encoded may be of any forms e.g., local addresses, addresses in private networks or Internet addresses, however, for sake of illustration, the examples given in the following description are based on URL type of
15 addresses. The address from which the certificate can be accessed can be encoded either when the file is transmitted from a server to the user system or when it is locally saved or transmitted to another system.

Figure 1 illustrates an example of the algorithm used
20 to encode a certificate address. As shown on figure 1a, a first step consists in getting the primary filename of the file (box 100), i.e. the filename of the file, and the address or URL of the certificate that is required to check the integrity and to verify the signature of the file,
25 referred to as certificate address in the following description (box 105). Then, the certificate address is encoded (box 110) and merged with the primary filename of the file, using particular separators (box 115) before the file is renamed with the filename comprising the primary filename
30 and the encoded certificate address (box 120).

Figure 1b depicts an example of the encoding algorithm (box 110). A variable *i* is set to zero (box 125) and the *i*th character is extracted from the certificate address string (box 130). A test is performed to determine whether the extracted character is valid or otherwise forbidden by filename syntax rules imposed by the file system of the user's device (box 135). If the extracted character is a filename valid character, variable *i* is incremented (box 150) and a test is performed to determine if variable *i* has reached its maximum value that is, if all characters of the certificate address string have been processed (box 155). If variable *i* has not reached its maximum value, the last four steps of the algorithm are repeated (boxes 130 to 155). Else, if variable *i* has reached its maximum value, the process is stopped. If the character extracted from the certificate address string is forbidden by the filename syntax rules, a corresponding valid character, or group of characters, is selected in lexicography table 145 and this selected character, or group of characters, replaces the forbidden one (box 140). Then variable *i* is incremented and the same test described before is performed to determine if variable *i* has reached its maximum value.

As an illustration of the algorithm described above, consider the case of a file based on Microsoft Word format (Word is a Trademark of Microsoft Corporation) named "Berry.doc", that a user would like to send to someone else as an e-mail attachment, using to this purpose a lexicography table to encode the certificate address string into the filename, wherein

30 ":" is associated to ".."
 "/" is associated to "("

To check the integrity and to verify the signature of this document file, it is required to use the certificate corresponding to the private key that has been used to sign this file. For sake of illustration one can considered that
5 this certificate can be downloaded from the following URL:

`http://www.certauth.com/certificates/Lewis Carroll.cer`

When the originator of the document "Berry.doc" signs the document, an option such as "Copy path to file" can be selected to encode the URL of the certificate repository
10 wherein the certificate required to check the integrity or to verify the signature of the document can be accessed.

The filename is modified according to the algorithm illustrated on figure 1. Firstly, by using the previous lexicography table, the certificate repository URL is
15 encoded as follows :

`http..((www.certauth.com(certificates(Lewis Carroll.cer`

Then, the encoded URL is merged with the filename. In this example, the encoded URL is enclosed in parenthesis that are used as separators. The encoded URL is inserted
20 front of the extension dot of the primary filename as follows :

`Berry(http..((www.certauth.com(certificates(Lewis Carroll.cer).doc`

and the file is renamed using this modified filename.

It must be noticed that, for sake of illustration, this
25 encoding algorithm is purposely very simple. A preferred one would consist in replacing a sequence of forbidden characters by a single one e.g., replacing "//:" by "(" . Likewise,

some sets of characters may be replaced by more compact codes e.g., replacing "http://" by "H!".

Figure 1c depicts an e-mail 160 wherein the filename 165 of the attached file 170 has been modified to embed the URL of the certificate address according to the previous algorithm.

When the attachment of the above mentioned e-mail is selected to be processed by the receiver, a test is performed to determine whether or not the user requests an integrity check or a signature verification so as to determine whether or not the certificate address must be extracted from the filename and decoded.

Using the same table of lexicographic transformations as the one that has been used by the sender of the file to encode the certificate address, the certificate address or URL is extracted and decoded from the filename. To that end, certain symbols or groups of symbols of the "encoded URL" are replaced by symbols or characters that are compatible with URL conventions on Internet, as mentioned above, to get the decoded and valid URL. Using the same example as before, the decoded certificate address is,

http://www.certauth.com/certificates/Lewis_Carroll.cer

Certificates are stored in a database of a certification authority server and, possibly, locally in the certificate's owner device. Each certificate comprises at least a public key that can be accessed by third parties to check the validity or to verify the signature of a signed file. The public key of a certificate corresponds to a private key that is known only by the certificate's owner and by the

certification authority, this private key being used to sign files. In a preferred embodiment, the certificates also comprise additional information such as the owner's name, the certificate's validity period and the signature
5 algorithm as mentioned above. It must be clear that a private key is only known by the certificate's owner and by the certification authority while all the other information relative to the private key and organized as a certificate is public and can be accessed by any third party knowing the
10 certificate address or URL.

Figure 2, comprising figures 2a and 2b, illustrates an example of the algorithm that is used to sign an electronic document and of the algorithm that is used to check the integrity and/or to verify the signature of a signed file,
15 respectively.

If the sender has not already a certificate issued by certification authority, he/she must apply for the certification authority to issue it. This must be made one time for a validity period since a certificate has a validity period.
20 Thus, the private key associated to a certificate issued by the certification authority can be used by the sender to sign all documents during the certificate validity period. To get a certificate the sender sends a request to the certification authority (step 200) with required information
25 such as sender's name. After having assigned a pair of private and public keys, the certification authority creates a certificate and transmits the private key as well as the certificate address to the user having sent the request, using a secure connection. The private key and the certificate
30 address are preferably stored locally on the user's device however, this information can be stored on a secure

server of the certification authority or on personal data storage means, such as a smart card.

After having selected the file to sign and once having received or recovered the required private key and the
5 associated certificate address (step 210), the user signs the file (step 215). To that purpose, a standard certification algorithm is used to compute a signature based on the file to be signed and the private key e.g. Message-Digest-5 (MD5) with RSA or SHA hashing algorithm with RSA.

10 In a preferred embodiment, the signature is appended to the document as illustrated on figure 4 wherein the signature (410) is located at the beginning of the file (400) and separated from the content of the document (405) by tags "BEGIN SIGNATURE" and "END SIGNATURE". Then, the address or
15 URL of the server wherein the public key that is required to check the integrity or to verify the signature of the file is encoded in the filename (step 220) as described by reference to figure 1. As mentioned above, the address or URL wherein this public key is stored is preferably provided by
20 the certification authority when issuing the certificate however, it can be transmitted to the sender, upon request, each time he/she signs a document. Therefore, at the end of the algorithm depicted on figure 2a, the resulting file is signed and contains a link to a server wherein a certificate
25 may be recovered to check the integrity of the resulting file or to verify the embedded signature.

Figure 2b illustrates an example of the algorithm that can be used to check the integrity or to verify the signature of a signed file encoded according to the invention.
30 The first step consists in decoding the filename of the

file, as described above, to retrieve the address or URL wherein the certificate that is required to check the integrity or to verify the signature is stored (step 225). Then, using this decoded address or URL, the user can access the certificate from a server, preferably controlled by a certification authority (steps 230 and 235), without opening the file. At this stage, the user can access information related to the certificate, such as the name of the person to whom the certificate has been delivered, the validity period of the certificate and the signature algorithm. Therefore, the user is able to check the certificate to determine whether or not the owner of the certificate is the one he/she expects to be (steps 240 and 245). Then, using the public key of the certificate, it is possible to authenticate the file i.e., to check the integrity of the file and/or to verify the signature (steps 250 and 255), by using a standard authentication algorithm. As suggested by dotted lines, the user can authenticate the signed file without checking the certificate. Naturally the certificate can contain information relative to the authentication algorithm that could be used, or must be used, depending upon the certification authority policy. Still in a preferred embodiment, the certification authority can provide the user means to download an authentication applet when he/she accesses the certificate so as to check the integrity and verify the signature of the file.

Figure 3 depicts an example of the environment wherein the invention can be implemented. For sake of illustration the main steps of the algorithms described on figures 2a and 2b are illustrated with referenced arrows. As described above, a user (300) who has no certificate and who wants to sign a file must access a certification authority server

(310) through a network (305) e.g., Internet. Certificates generated by the certification authority (320) are locally stored in a certificate database (315) of the certification authority server (310). Likewise, when a user (325) having a signed file e.g., received as an e-mail attachment, wants to check its integrity and to verify the signature, he/she accesses through a network (305) the public key of the certificate which address or URL is encoded in the filename of the signed file to check.

10 Figure 4 shows a signed file (400) comprising the document (405) and a signature (410) that can be used to check the file integrity and to verify the identity of the document's author. The address or URL of the certification authority server wherein the certificate corresponding to
15 the private key used to sign the file is encoded and stored in the filename (415).

Figure 5, comprising figures 5a to 5f, illustrates an example of the user's interface when using the invention. Figures 5a to 5d depict an example of certificate panel
20 while figures 5e and 5f show how a certificate address or URL can be linked to a file.

The certificate panel illustrated on figures 5a to 5d comprises 4 tabs depicted on each of these figures, respectively, these tabs comprising information relative to:

- 25 - general tab:
- owner of the certificate,
 - certification authority having delivered the certificate, and,
 - validity period,

- detail tab:
 - version identifying the certificate format,
 - serial number (unique within the certification authority),
 - 5 • signature algorithm (identifying the issuer's hash algorithm and digital signature algorithm used to sign the certificate),
 - issuer name (the name of the certification authority),
 - 10 • the beginning of the validity period,
 - the end of the validity period,
 - subject name (the name of the user to whom the certificate is issued),
 - subject's public key field (including Algorithm name and the Public Key of the subject),
 - 15 • extensions, and,
 - signature of the certification authority,
 - certification path (address or URL of the certificate on the certification authority server), and,
 - 20 - download SW (comprises links software applications or applets that are adapted, for example, to check the validity of a file or to verify a signature).

Most of these fields are completed by the certification authority after having received a request for a certificate and an identifier or subject name. The private and public keys are computed according to standard algorithms.

Figures 5e and 5f depict an example of the interface that can be used to encode a certificate address or URL into the filename of a file. After having selected a file in the file manager, the user can click on the right button of the

mouse to display a pop-up menu comprising a "Paste path to
file" option. Then, the path previously memorized in the
clipboard or selected by other means is encoded in the
filename of the file according to the method described by
5 reference to figure 1.

Naturally, in order to satisfy local and specific
requirements, a person skilled in the art may apply to the
solution described above many modifications and alterations
all of which, however, are included within the scope of
10 protection of the invention as defined by the following
claims.

Claims:

1. A method for encoding in the filename of a signed file, an address from which the certificate required to authenticate said signed file can be accessed, said method comprising the steps of,
- 5 - encoding said address from which the certificate required to authenticate said signed file can be accessed (step 110);
- 10 - merging said filename and said encoded address in a new filename (step 115); and,
- renaming said signed file with said new filename (step 120),
- wherein said filename and said encoded addresses are separated by a control character.
- 15 2. The method of claim 1 wherein said step of encoding the address from which the certificate required to authenticate said signed file can be accessed comprises the steps of :
- 20 - analyzing the address from which the certificate required to authenticate said signed file can be accessed to detect predetermined characters (step 135); and,
- replacing said predetermined characters by associated characters (step 140),
- said predetermined and associated characters being stored in a lexicography table (145).
- 25 3. The method of either claim 1 or claim 2 wherein the encoded address is merged with said filename, by inserting

the encoded address in front of the extension dot of said filename.

4. The method of any one of claims 1 to 3 further comprising a compression step to reduce the size of the encoded address.

5. A method for authenticating a signed file having a filename wherein an address from which the certificate required to authenticate this signed file can be accessed is encoded, said method comprising the steps of,

10 - extracting said encoded address;

 - decoding said encoded address (step 225);

 - accessing said certificate required to authenticate said signed file using said decoded address (steps 230 and 235),

15 - authenticating said signed file using said accessed certificate (steps 240 to 255).

6. The method of claim 5 further comprising the step of checking said certificate before authenticating said signed file.

20 7. The method of either claim 5 or claim 6 wherein said step of accessing said certificate required to authenticate said signed file further comprises the step of,

 - downloading an authentication function according to information contained in said certificate.

8. The method of any one of claims 5 to 7 wherein said step of authenticating said signed file comprises at least one of the steps consisting in checking the integrity of said signed file and in verifying the signature.
- 5 9. The method of any one of claims 5 to 8 further comprising the step of extracting a public key from said certificate, said public key being used to authenticate said signed file.
- 10 10. The method of any one of claims 1 to 9 wherein said certificate is stored in a server of an authentication authority.
11. An apparatus comprising means adapted for carrying out each step of the method according to any one of the claims 1 to 10.
- 15 12. A computer-like readable medium comprising instructions for carrying out each step of the method according to any one of the claims 1 to 10.

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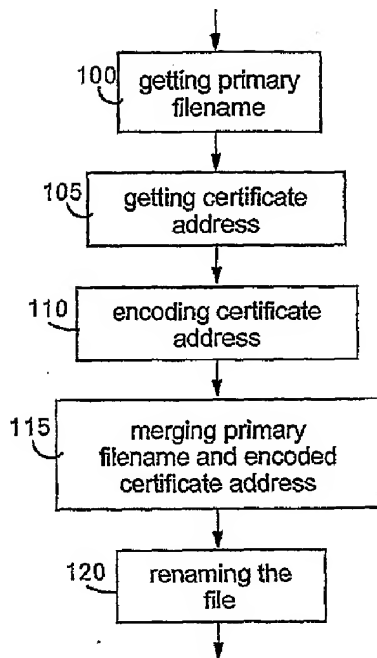


Figure 1a

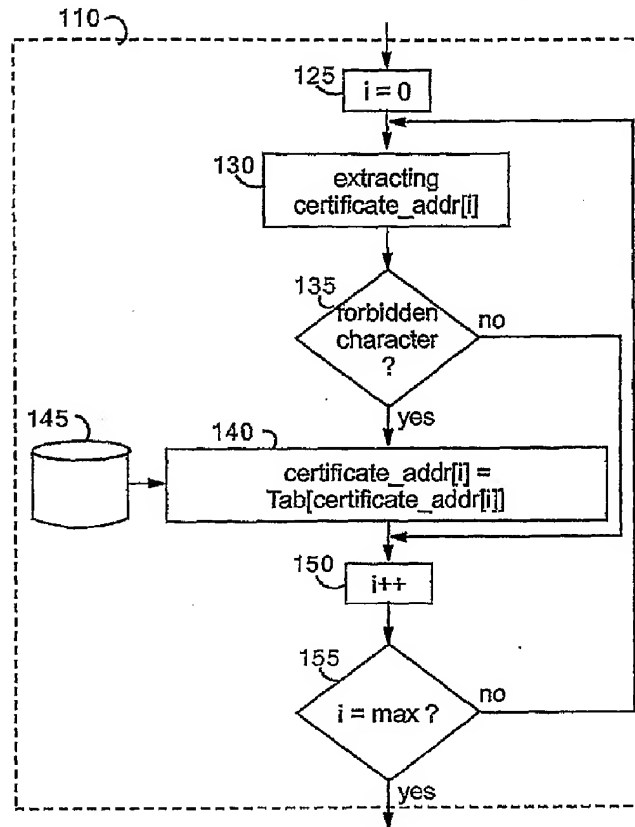


Figure 1b

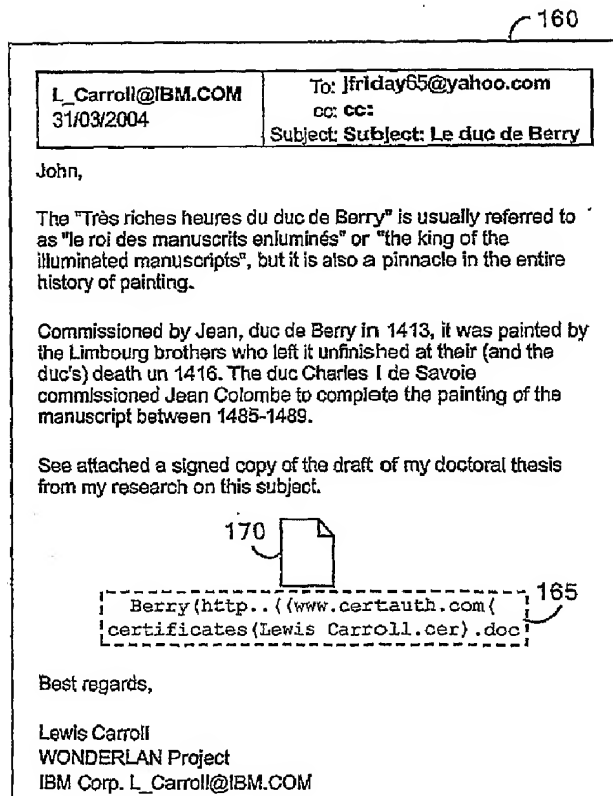


Figure 1c

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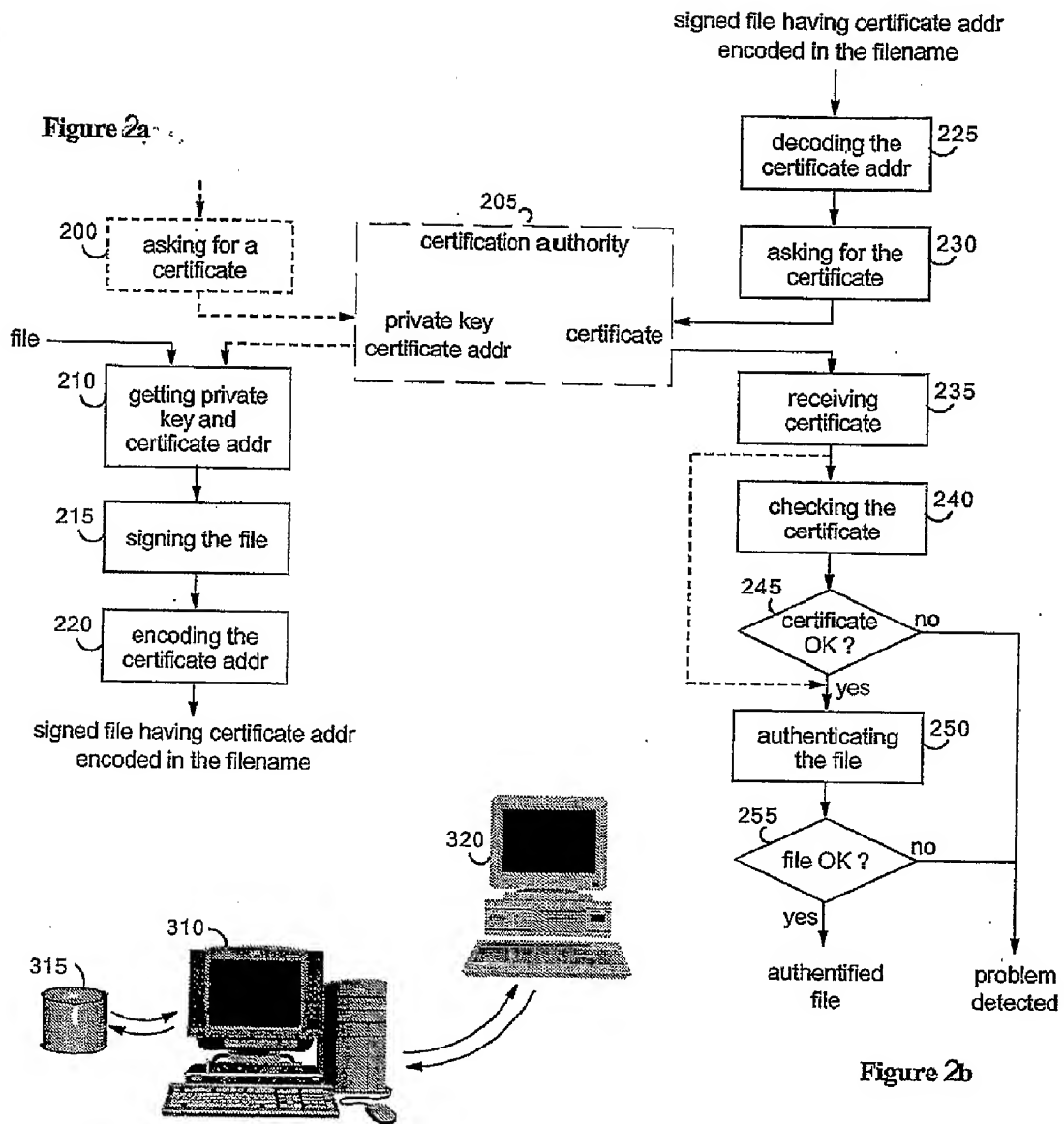
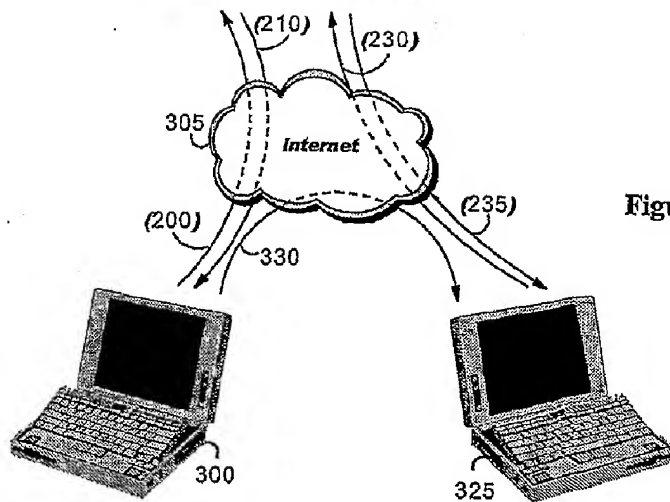


Figure 2b






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LES TRÈS RICHES HEURES DU DUC DE BERRY Lewis Carroll, WONDERLAND Project, IBM Corp.	
WHAT IS THE TRÈS RICHES HEURES? The Tres Riches Heures is the classic example of a medieval book of hours. This was a collection of the text for each liturgical hour of the day - hence the name - which often included other, supplementary, texts. Calendars, prayers, psalms and masses for certain holy days were commonly included.	
The pictures in this directory are from the calendar section of the Tres Riches Heures. This was painted some time between 1412 and 1416 and is arguably the most beautiful part of the manuscript; it is certainly the best known, being one of the great art treasures of France. In terms of historical and cultural importance, it is certainly equal to more famous works such as the Mona Lisa, marking the pinnacle of the art of manuscript illumination.	
WHO PAINTED THE TRÈS RICHES HEURES? The Tres Riches Heures was painted by the Limbourg brothers, Paul, Hermann and Jean. They came from Nimwegen in what is now Flanders but were generally referred to as Germans. Very little is known about them; they are believed to have been born in the late 1370s or 1380s and were born into an artistic family, their father being a wood sculptor and their uncle being an artist working variously for the French Queen and for the Duc de Bourgogne.	
They seem to have followed in their uncle's footsteps and by 1402 had entered into the service of the Duc de Bourgogne as artists. By 1408 they had entered the service of Jean, Duc de Berry, one of the most notable (and richest!) art lovers in France. They are known to have executed several other pieces of work apart from the Tres Riches Heures but most	

Figure 4

Berry (http..(www.certauth.com(certificates(Lewis Carroll.cer).doc

Figure 5a

Certificate	
General	Details
Certificate Path Download SW	
 Certificate Information	
This certificate is intended to: <ul style="list-style-type: none"> • Ensure the identity of this certificate owner 	
<hr/> Issued to: Lewis Carroll	
Issued by: Certification Authority	
Valid from: 16/06/2003 to 16/07/2005	
<div style="text-align: right;">   </div>	

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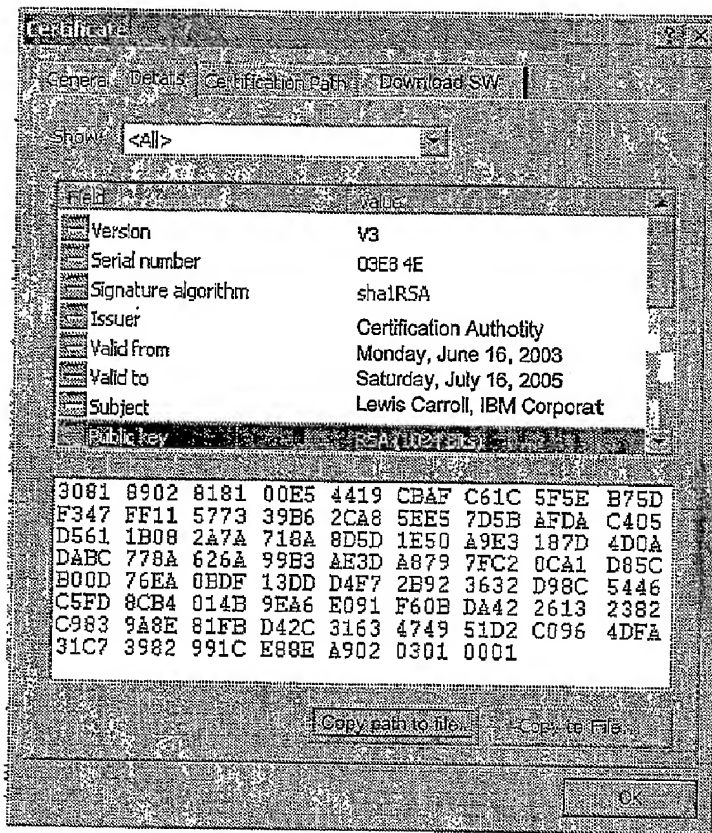
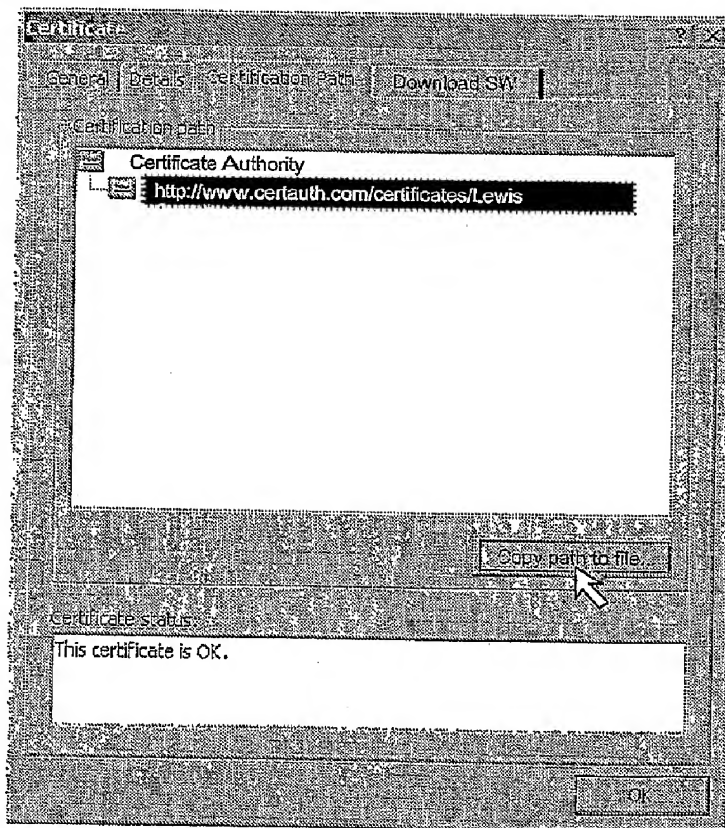


Figure 5b

Figure 5c



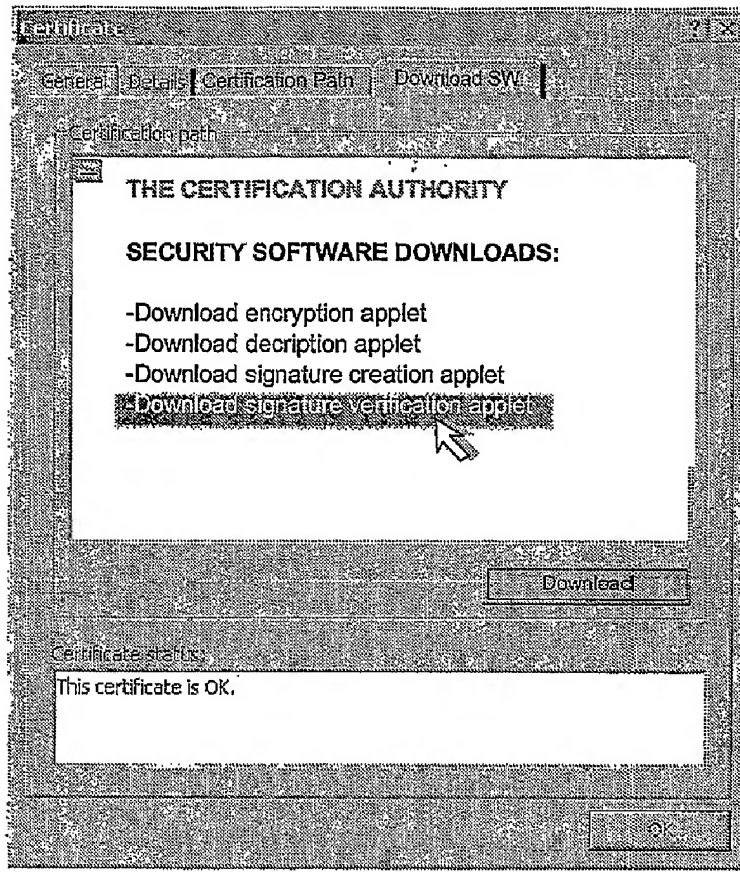


Figure 5d

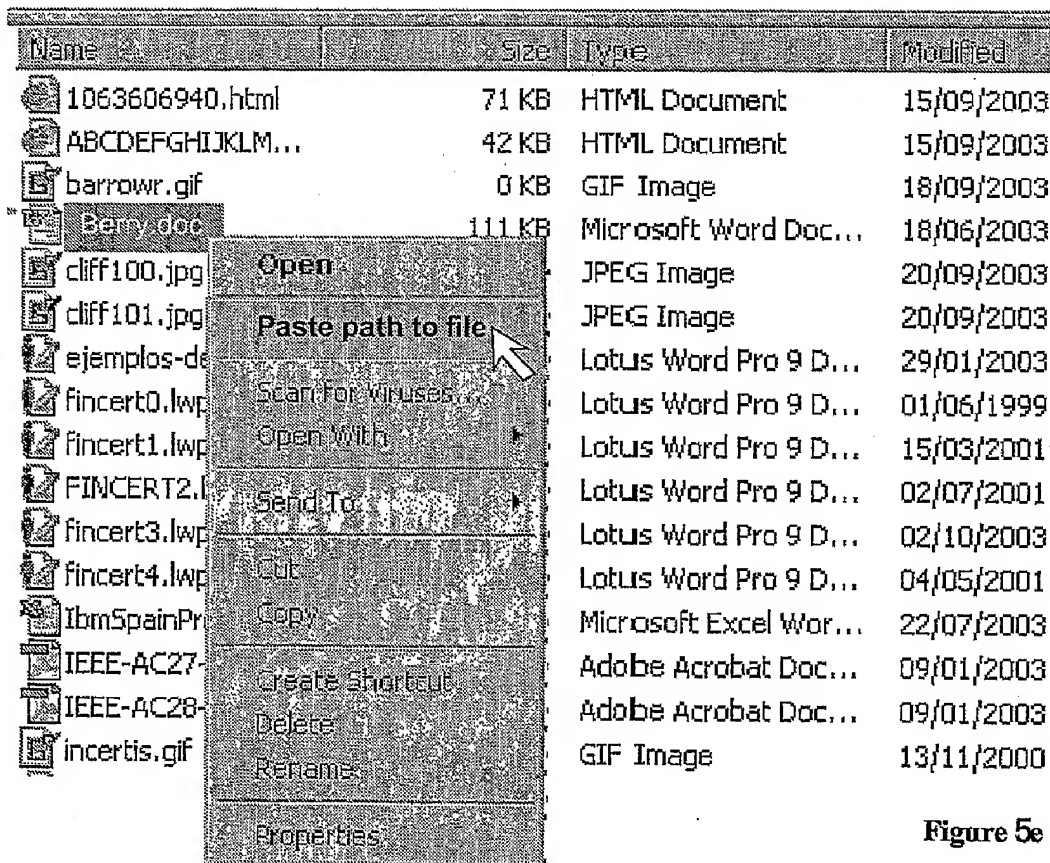


Figure 5e

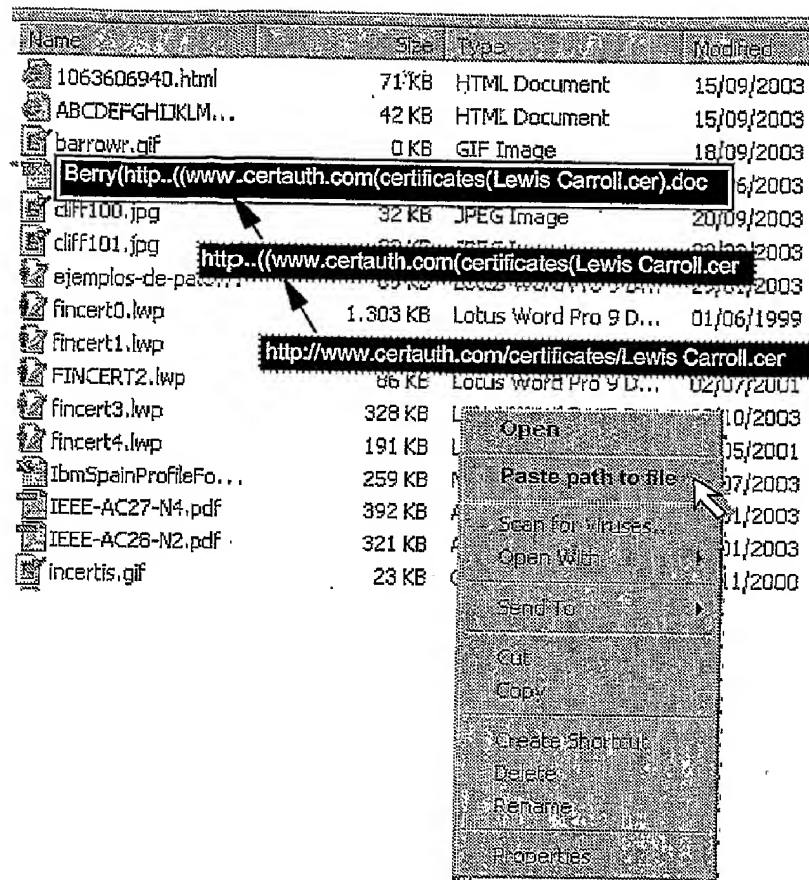


Figure 5f

INTERNATIONAL SEARCH REPORT

corrected version

International Application No
PCT/EP2005/050925A. CLASSIFICATION OF SUBJECT MATTER
IPC 7 G06F1/00

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC 7 G06F

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

EPO-Internal

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X	HELPPFILE OF VERSION 2.9 OF MP3 TAG CLINIC, [Online] 10 February 2004 (2004-02-10), XP002334789 Retrieved from the Internet: URL: http://web.archive.org/web/20040210003723/http://www.kevesoft.com/files/clinic29.zip [retrieved on 2005-07-04]	1-4, 11, 12
A	"Rename Page", "Illegal Character Editor", "Main Editor", "Options Panel", "URL Editor", "Metadata Attribute List / Cross Reference", from the Index of the HelpFile of version 2.9 of Mp3 Tag Clinic. the whole document	5-10
X	US 2001/020272 A1 (LE PENNEC JEAN-FRANCOIS ET AL) 6 September 2001 (2001-09-06) paragraph [0052] - paragraph [0076] paragraph [0169] - paragraph [0194] -/-	5-12

☒ Further documents are listed in the continuation of box C.☒ Patent family members are listed in annex.

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"O" document referring to an oral disclosure, use, exhibition or other means

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"Y" document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art.

"Z" document member of the same patent family

Date of the actual completion of the international search

5 July 2005

Date of mailing of the international search report

18. 08. 2005

Name and mailing address of the ISA

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Authorized officer

Alecú, M

INTERNATIONAL SEARCH REPORT

corrected version

International Application No

PCT/EP2005/050925

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X	& EP 0 813 132 A (INTERNATIONAL BUSINESS MACHINES CORPORATION) 17 December 1997 (1997-12-17) page 2, line 29 - line 45 -----	5-12
A	SUEL T ET AL: "Compressing the graph structure of the Web" DATA COMPRESSION CONFERENCE, PROCEEDINGS. DCC, IEEE COMPUTER SOCIETY PRESS, LOS ALAMITOS, CA, US, 2001, pages 213-222, XP002285665 page 217 - page 219 -----	4
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A	ASHMAN H: "Theory and Practice of Large-Scale Hypermedia Systems" 1997, , XP002285663 page 45, paragraph 3.3 - page 47 -----	
A	"X.500 Object Identifiers allocated by OpenFortress" [Online] 19 June 2002 (2002-06-19), XP002334790. Retrieved from the Internet: URL: http://web.archive.org/web/20020619230645/http://openfortress.nl/doc/OID/1.3.6.1.4.1.10471.html > [retrieved on 2005-07-04] page 1, paragraph 0.1 -----	1

INTERNATIONAL SEARCH REPORT

Information on patent family members

corrected version

International Application No

PCT/EP2005/050925

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US 2001020272 A1	06-09-2001	NONE	
EP 0813132 A	17-12-1997	US 5825877 A	20-10-1998
		DE 69732323 D1	03-03-2005
		EP 0813132 A2	17-12-1997
		JP 10083310 A	31-03-1998
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